	ON PAGE	AAA QV:11/CC	<del></del>	Form Approved OMB No. 0704-		
AD-A216 670		16. RESTRICTIVE	MAKKINGS			
		3. DISTRIBUTION/AVAILABILITY OF REPORT  Approved for public nations; distribution unlimited.  5. MONITORING ORGANIZATION REPORT NUMBER(S)  APOSR - TR - 89 - 0570				
6a. NAME OF PERFORMING ORGANIZ	ATION 6b. OFFICE SYMBOL (If applicable)	1_				
Stanford Houses	AFOSR	AFOSR				
Standord Elningers to			7b. ADDRESS (City, State, and ZIP Code)			
660 Augus 110 Way Stanford, CA 9430		Building Bolling,	410 AFB DC 20	332–6448		
a. NAME OF FUNDING / SPONSORING 8b. OFFICE SYMBOORGANIZATION (If applicable)		9. PROCUREMENT INSTRUMENT IDENTIFICATION NUMBER				
AFOSR NM		AFOSR- 87-0149				
Building 410 Building, AFB DC 20332-6448  11. TITLE (Include Security Classification)		PROGRAM ELEMENT NO.	PROJECT NO.	TASK NO	WORK I	
		61102F	2304	A:	1714	
Deductive Comput	rea Programmina				ELECTE JANO91	
12. PERSONAL AUTHOR(S)	4				JANOS	
PROFESSOR MANNA  13a. TYPE OF REPORT  FROM 1 Mar 8 TO 28 Feb 88		14. DATE OF REPO	PRT (Year, Mon	th, Day.	90	
16. SUPPLEMENTARY NOTATION				<b>_</b>		
17. COSATI CODES	18. SUBJECT TERM	S (Continue on rever	e if necessary	and identify	by block number	
FIELD GROUP SUB-C	GROUP					

proof/theoretic deduction plays a crucial role in understanding logic programs. Moreover, for specific implementations of logic programming, e.g. PROLOG, the notion of deduction stategy is also important.

and en organic Conquery of the	-89 5 11 098°
20. DISTRIBUTION/AVAILABILITY OF ABSTRACT UNCLASSIFIED/UNLIMITED SAME AS RPT. DTIC USERS	21. ABSTRACT SECURITY CLASSIFICATION  OCCUPY
2a. NAME OF RESPONSIBLE INDIVIDUAL	22b. TELEPHONE (Include Area Code) 22c. OFFICE SYMBOL (202) 767-

S PAGE

# DEDUCTIVE COMPUTER PROGRAMMING

AFOSR-TR- 89-0570

Final Technical Report:
Department of the Air Force
Grant AFOSR 87-0149
Expiration Date: February 29, 1988



by
Zohar Manna, Professor
Computer Science Department
Stanford University
Stanford, California 94305

Acces	aien For	. •
	GRA&I	1
DTIC		j
	eunced []	
Justi	fication	-{
		1
77	About A and A	┥
	ibution/	4
Avai	lability Codes	1
	Avail and/or	1
Dist	Special	1
! }		İ
	] ]	ł
A-I		
	<i></i>	185 Y
	1	615
	<i>j</i>	41
	\	<b>_</b>

#### TECHNICAL SUMMARY

Our research concentrated on the following topics:

#### • Logic Programming Semantics: Techniques and Applications ([B1]-[B3])

It is generally agreed that providing a precise formal semantics for a programming language is helpful in fully understanding the language. This is especially true in the case of logic-programming-like languages for which the underlying logic provides a well-defined but insufficient semantic basis. Indeed, in addition to the usual model-theoretic semantics of the logic, proof-theoretic deduction plays a crucial role in understanding logic programs. Moreover, for specific implementations of logic programming, e.g. PROLOG, the notion of deduction stategy is also important.

We provided semantics for two types of logic programming languages and develop applications of these semantics. First, we propose a semantics of PROLOG programs that we use as the basis of a proof method for termination properties of PROLOG programs. Second, we turn to the temporal logic programming language TEMPLOG of Abadi and Manna, develop its declarative semantics, and then use this semantics to prove a completeness result for a fragment of temporal logic and to study TEMPLOG's expressiveness.

In our PROLOG semantics, a program is viewed as a function mapping a goal to a finite or infinite sequence of answer substitutions. The meaning of a program is then given by the least solution of a system of functional equations associated with the program. These equations are taken as axioms in a first-order theory in which various program properties, especially termination or non-termination properties, can be proved. The method extends to PROLOG programs with extra-logical features such as cut.

For TEMPLOG, we provide two equivalent formulations of the declarative semantics: in terms of a minimal temporal Herbrand model and in terms of a least fixpoint. Using the least fixpoint semantics, we are able to prove that TEMPLOG is a fragment of temporal logic that admits a complete proof system. This semantics also enables us to study TEMPLOG's expressiveness. For this, we focus on the propositional fragment of TEMPLOG and prove that the expressiveness of propositional TEMPLOG queries essentially corresponds to that of finite automata.

#### The TABLOG language and its implementation ([MMW])

Logic programming uses formal proofs as the computation paradigm. That is, a logic program is a theory, expressed in a given logic, that captures some properties of the real world. The execution of such a program is the proof of some theorem in this theory.

TABLOG is a new logic-programming language ([M][MMW1][MMW2]) based on quantifier-free first-order logic with equality, using the proof rules of the deductive-tableau theorem-proving method as the execution mechanism.

The main features of TABLOG are consequences of the use of full first-order logic. In particular, TABLOG incorporates all the standard connectives, not only implication and conjunction, but also equality, negation and equivalence. Programs are nonclausal: they do not need to be in Horn-clause form or any other normal form. Programs can compute relations (as in PROLOG) or functions (as in LISP), whichever is more appropriate; this improves the clarity and the efficiency of programs.

Terms are lazy-evaluated to make the use of functions more convenient. No cut annotation is required as the system can detect such optimizations dynamically.

Three deduction rules are used for the execution of the programs: nonclausal resolution (case analysis), equality replacement (repalcement of equal terms), and equivalence replacement (repalcement of equivalent subsentences).

We have developed of a compiler for TABLOG; this compiler will produce code for a virtual TABLOG machine, similar to the Warren abstract machine. This compiler, written in TABLOG itself, will support a new syntax, which includes types and an elaborate notion of modules and generic modules. The virtual machine was implemented on a Sun workstation.

### • Logic: The Calculus of Computer Science ([MW])

The research papers in which we have presented the deductive approach to program synthesis has been addressed to the usual academic readers of the scholarly journals. In an effort to make this work accessible to a wider audience, including computer science undergraduates and programmers, we have developed a more elementary treatment in the form of a two-volume book, *The Logical Basis for Computer Programming*, Addison-Wesley (Manna and Waldinger [85c]).

This book requires no computer programming and no mathematics other than an intuitive understanding of sets, relations, functions, and numbers; the level of exposition is elementary. Nevertheless, the text presents some novel research results, including

- theories of strings, trees, lists, finite sets and bags, which are particularly well suited to theorem-proving and program-synthesis applications;
- formalizations of parsing, infinite sequences, expressions, substitutions, and unification;
- a nonclausal version of skolemization;
- a treatment of mathematical induction in the deductive-tableau framework.

## • Verification of Concurrent Programs ([MP])

We studied in detail the proof methodologies for verifying temporal properties of concurrent programs. Corresponding to the main classification of temporal properties into the classes of safety and liveness properties, appropriate proof principles were presented for each of the classes.

We developed proof principles for the establishment of safety properties. We showed that essentially there is only one such principle for safety proofs, the invariance principle, which is a generalization of the method of intermediate assertions. We also indicated special cases under which these assertions can be found algorithmically.

The proof principle that we developed for *liveness* properties is based on the notion of well-founded descent of ranking functions. However, because of the nondeterminancy inherent in concurrent computations, the well-founded principle must be modified in a way that is strongly dependent on the notion of fairness that is assumed in the computation. Consequently, three versions of the well-founded principle were presented, each corresponding to a different definition of fairness.

#### REFERENCES

Research papers and Ph.D. theses supported by this contract.

- \* indicates papers that are attached as part of this report.
- \*[B1] M. Baudinet, "Proving Termination Properties of PROLOG Programs: A Semantic Approach", Proceedings of the Third Annual Symposium on Logic in Computer Science., pp. 336-347, Edinburgh, Scotland, July 1988.
- \*[B2] M. Baudinet, "Temporal Logic Programming is Complete and Expressive", Proceedings of the Sixteenth ACM Symposium on Principles of Programming Languages, Austin, Texas. January 1989.
- [B3] M. Baudinet (supervised by Z. Manna), Logic Programming Semantics: Techniques and Applications, Ph.D. Thesis, Computer Science Dept., Stanford University (1989).
- [MP] Z. Manna and A. Pnueli, Temporal Specification and Verification of Concurrent Program. Textbook, to appear 1989.
- [MW] Z. Manna and R. Waldinger, Logical Basis for Computer Programming, Textbook, Addison-Wesley Pub., Volume 2: Deductive Systems (to appear, 1989).
- \*[MMW] E. Muller, Z. Manna, and R. Waldinger, "The TABLOG Language," Draft, 1988.